

PACE

Simple Multi-hop Scheduling for Single-radio 802.11-based Stub Wireless Mesh Networks



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Outline

Introduction

IEEE 802.11-based Stub WMN

Stub WMN Major Problems

PACE

Simulation results

Conclusions and future work

Introduction

facebook



Internet playing central role



twitter

Google

You Tube
Broadcast Yourself™

Ubiquitous wireless connectivity to the Internet

Increasing number
of mobile devices

Large number of
Access Points

Complex
deployment

Limited wired
connectivity

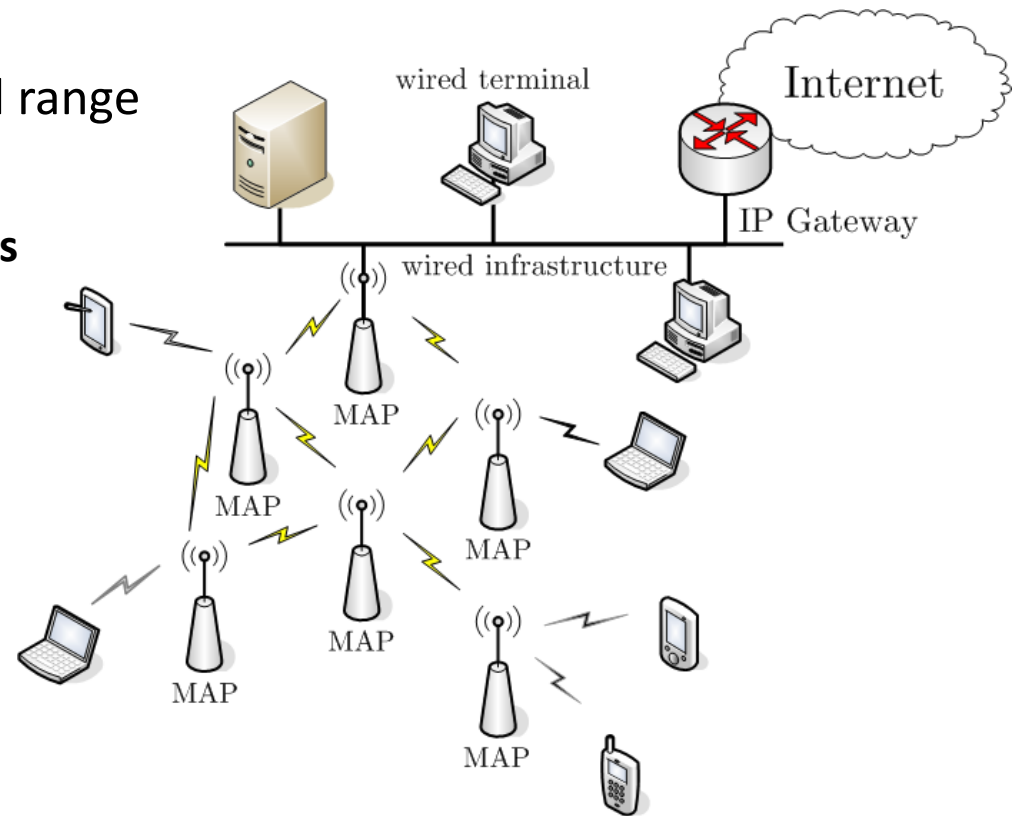


IEEE 802.11-based Stub Wireless Mesh Networks

IEEE 802.11-based Stub WMN

Advantages

- Compliant with legacy systems
 - **Infrastructure**
 - **STAs**
- Infrastructure extended range
- Low cost deployment
 - **Free radio frequencies**

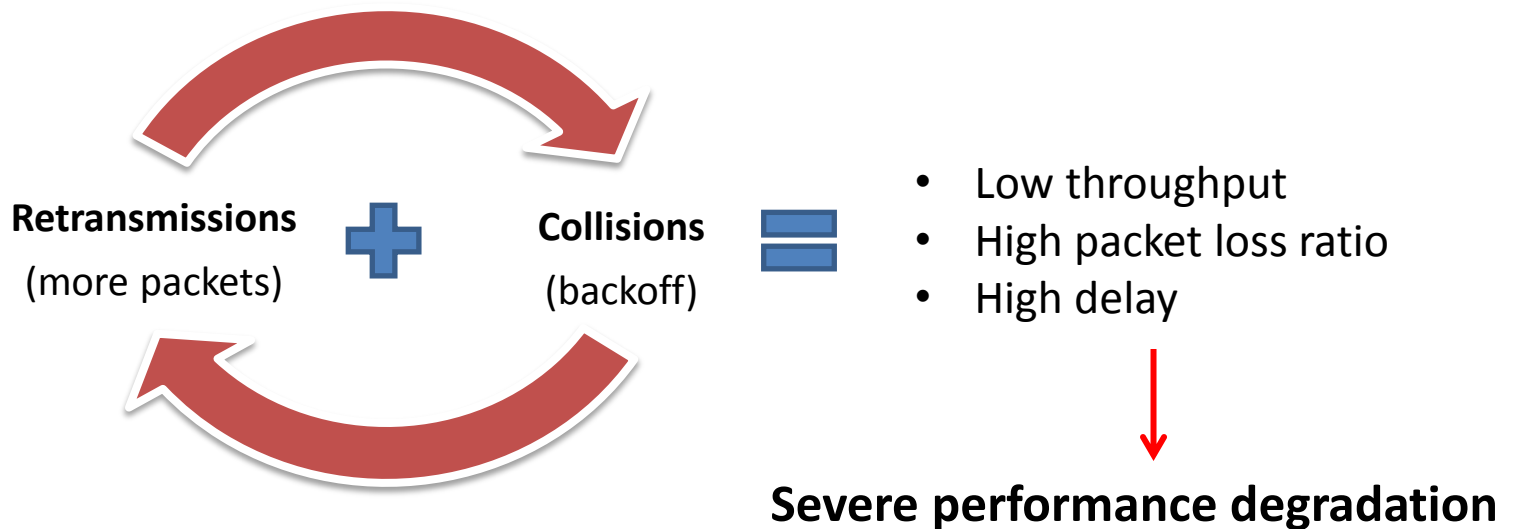


802.11-based Stub WMN – Major Problems

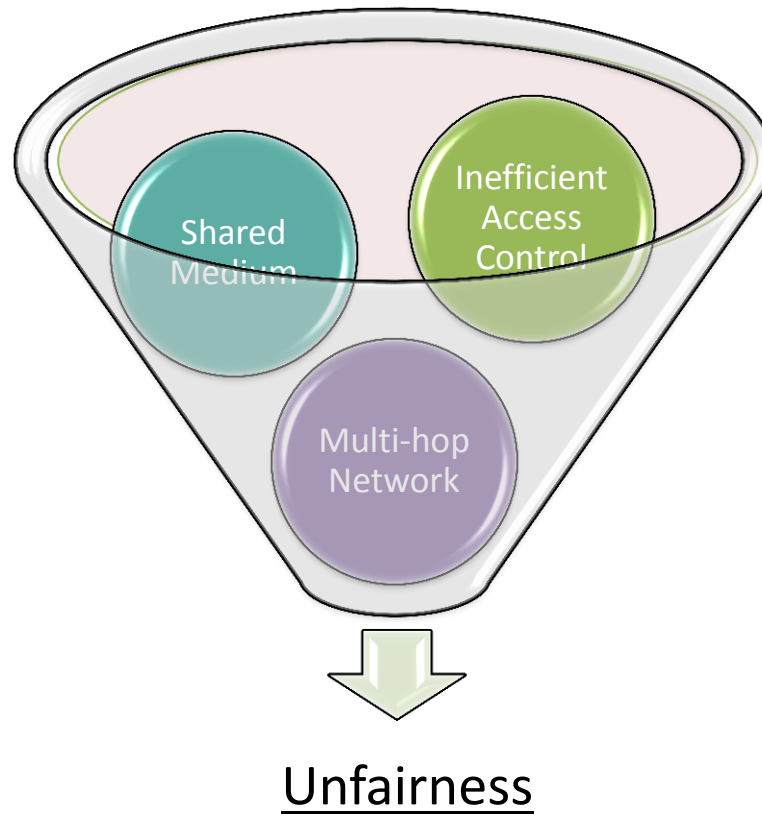
Inefficiency

CSMA/CA designed for single-hop networks

- Unable to avoid frame collisions in multi-hop networks
 - Inter/intra-flow interference (spatial contention over multiple hops)
 - Hidden/exposed node problems worsened
 - RTS/CTS harmful



802.11-based Stub WMN – Major Problems



Nodes closer to root take control of the medium → **Node starvation**

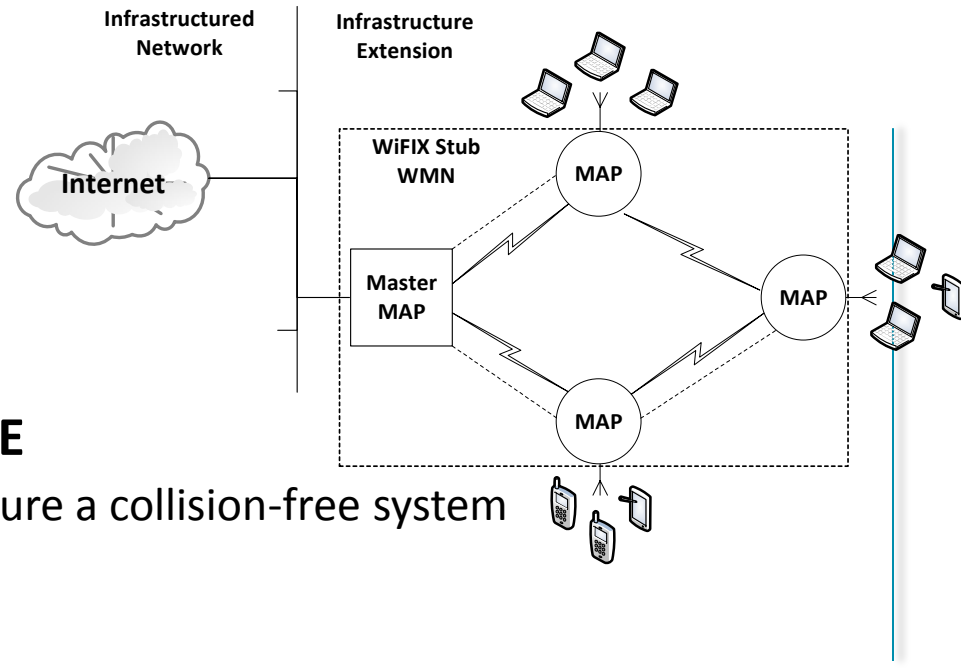
WMN – state of the art approaches

- Examples
 - RTS/CTS modifications & improvements
 - Clique techniques
 - Radio strategies
 - Spatial reuse / scheduling (NP-hard optimization problem)
- High complexity
 - Time synchronization (fixed size time slots)
 - Modifications to MAC layer
 - Higher overhead
 - Higher equipment costs
- Restrictive
 - Solve specific problems or scenarios

Simple, single-radio, 802.11-based solution is needed!

WiFIX – Wi-Fi network Infrastructure eXtension

- Tree-based routing solution for Stub WMN
- Runs on top of 802.11 MAC
- Active tree topology creation/maintenance
- Layer-2 routing based on IEEE 802.1D bridges
- Virtual Ethernet links between neighbours



Upgraded to accommodate PACE

- Routing technique updated to ensure a collision-free system
- Tweaked to improve performance


PACE – Simple Multi-hop Scheduling Mechanism



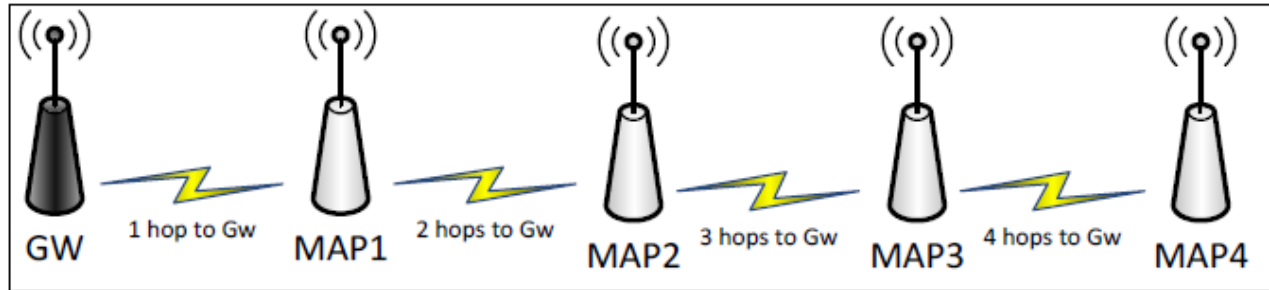
PACE – Major characteristics

- Data plane collision free operation
 - Master controls multi-hop scheduling
 - Single packet travels in Stub WMN
 - Suitable for high offered loads
- Advantages
 - Fairness between all MAPs
 - Network capacity equally divided (no spatial bias, unless enforced)
 - No starvation in MAPs far from the gateway
 - Predictable capacity
 - $C = C_{\text{Wi-Fi link}} / \text{AvgHopCount}$
 - QoS easily implementable
 - Managed by the master
- Main disadvantages
 - Exposed nodes as number of hops increases -> no spatial reuse
 - Less efficient for low offered loads

PACE – Implementation and Testing

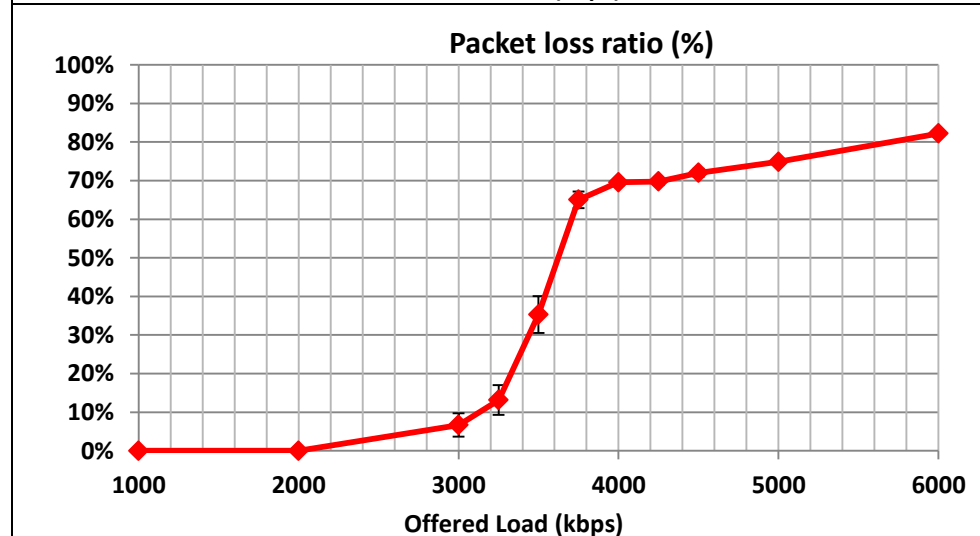
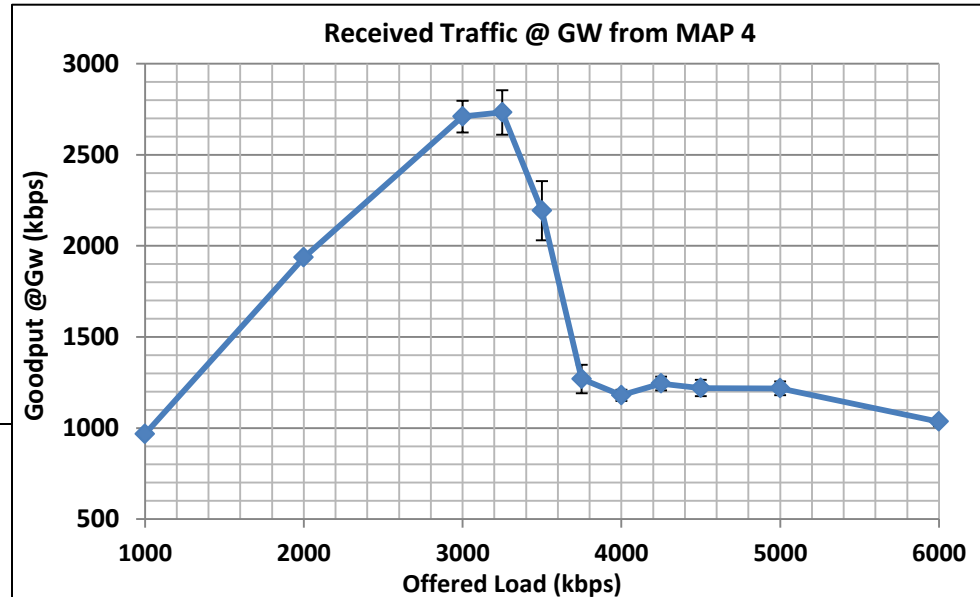
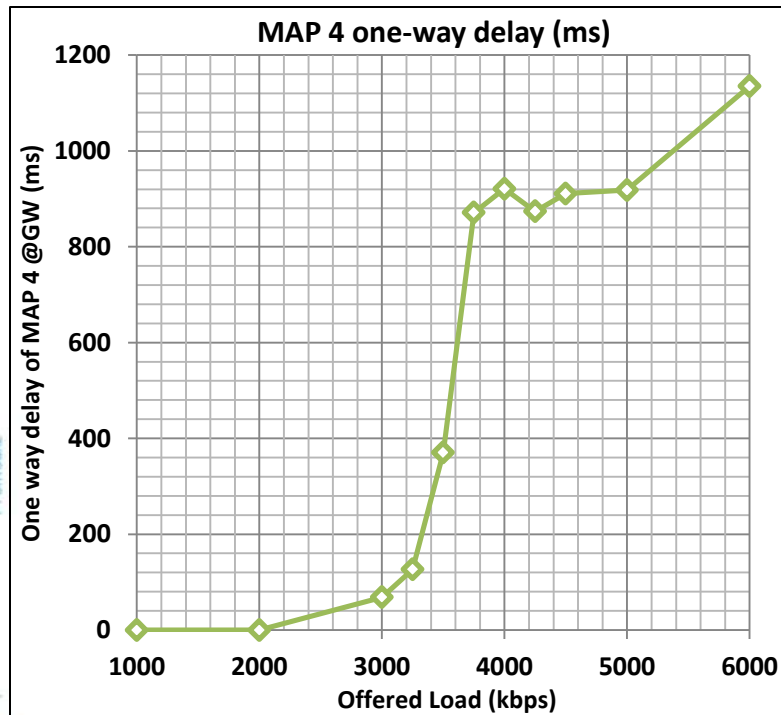
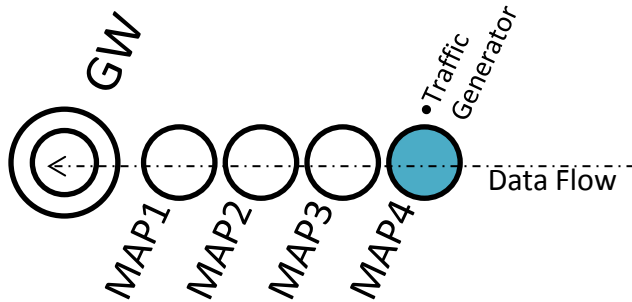
- Implemented in 
- Steps:
 1. WiFIX full implementation (with the upgrade required by PACE)
 2. PACE integration
 3. Simulation
 - i. **Fixed topology**
 - a. Chain with a gateway and 4 MAPS
 - b. 7 use cases studied
 - ii. **Random topologies**
 - a. MAPs arranged in a radial topology with the gateway at the centre
 - b. 1 use case studied
 4. Results gathered in multiple use cases
 - i. **Native CSMA/CA**
 - ii. **PACE**

Scenario 1 – Fixed Topology | 7 Use cases

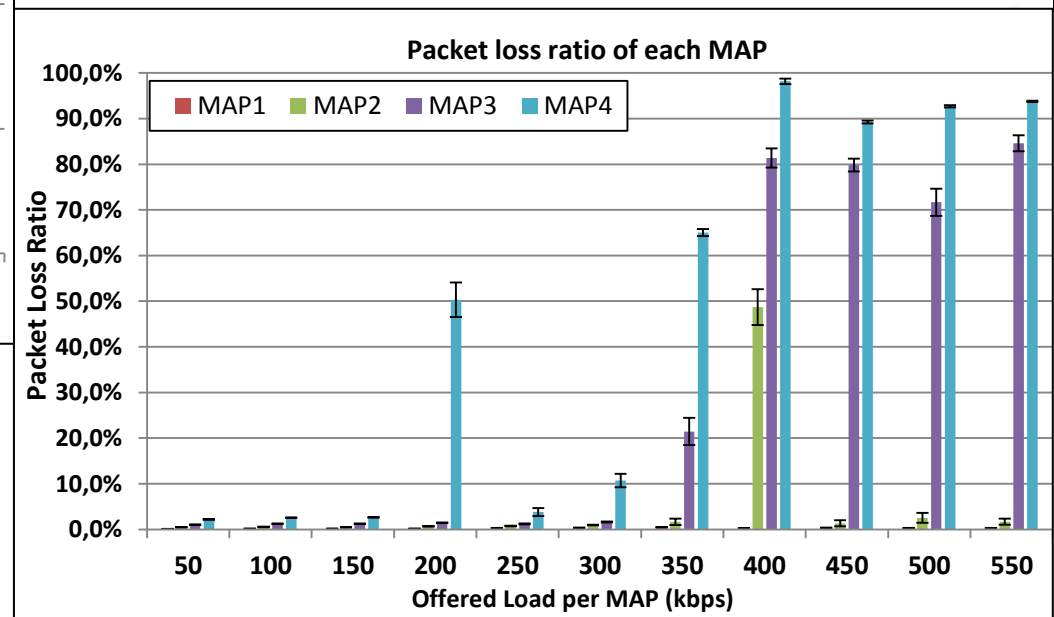
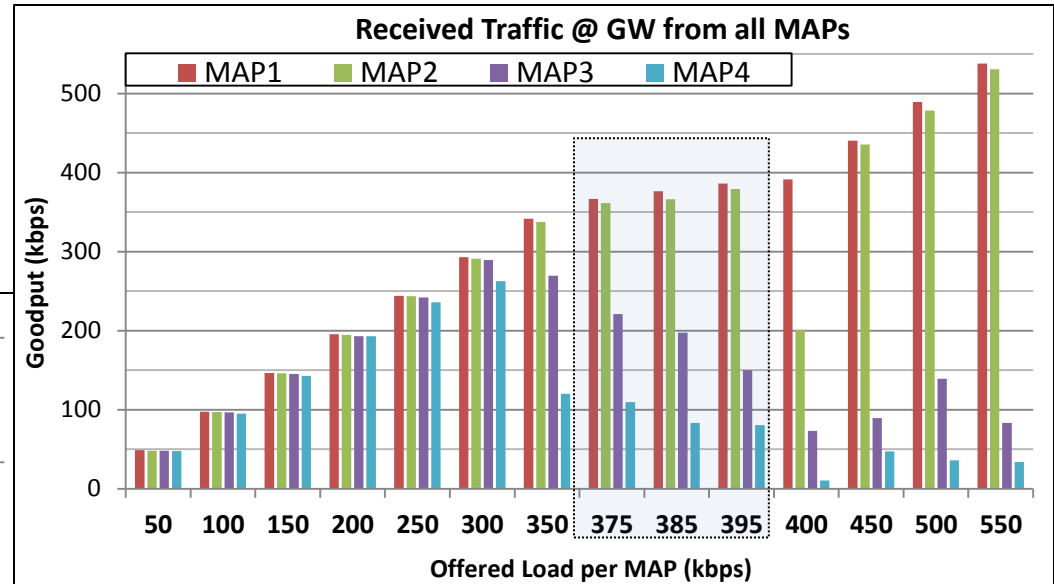
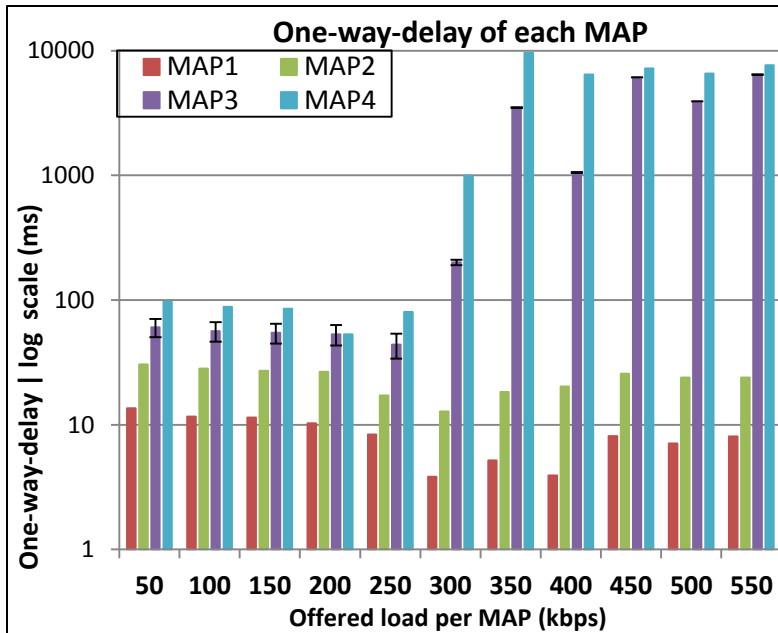
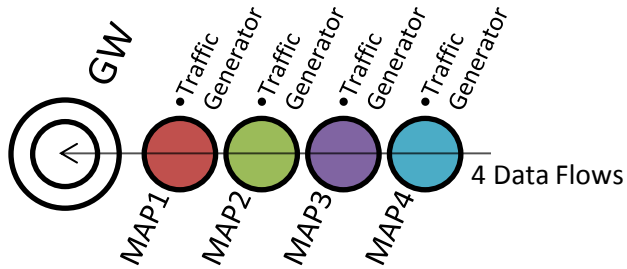


- One Gateway and four MAPs in a chain topology
- Size of data frames: 2000 bytes
- RTS/CTS disabled
- Results are expressed as Goodput and Delay

Results – Case 1 | CSMA/CA only | Single flow through a chain of nodes



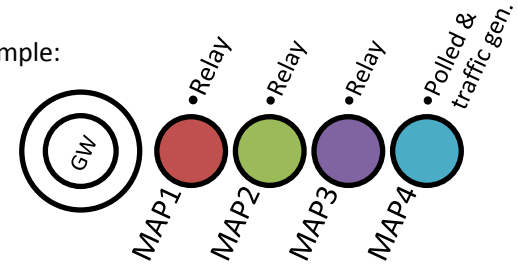
Results – Case 2 | CSMA/CA only | Multiple flows through a chain of nodes



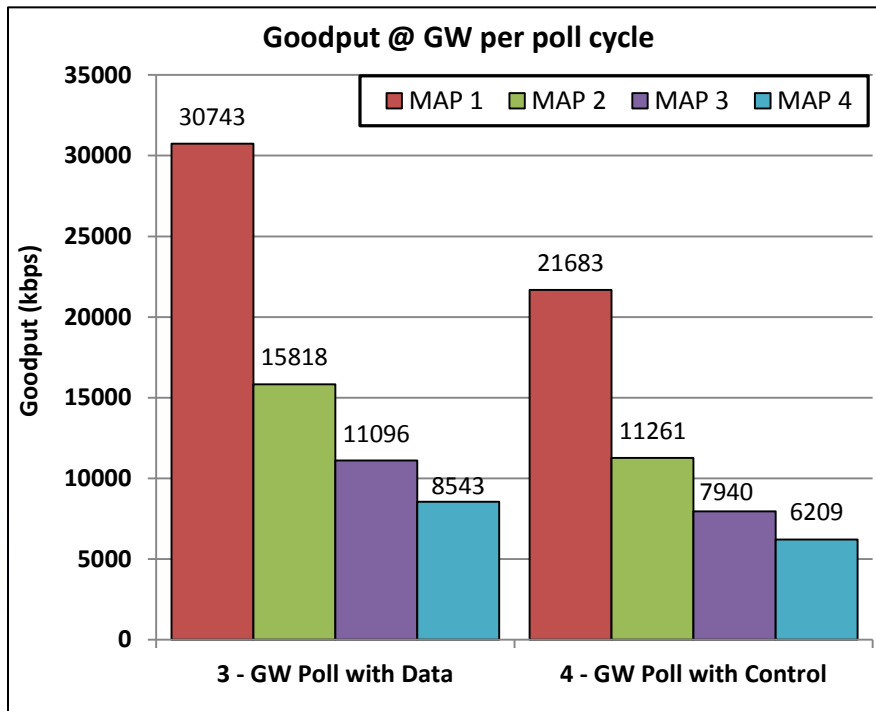
Results – Cases 3 and 4 | PACE | Direct MAP polling

- Gateway polls directly a specific MAP repeatedly
- Polled MAPs are active: Traffic generator
- In Case 3 the poll signal is implicit (Data)
- In Case 4 the poll signal is explicit (Control)

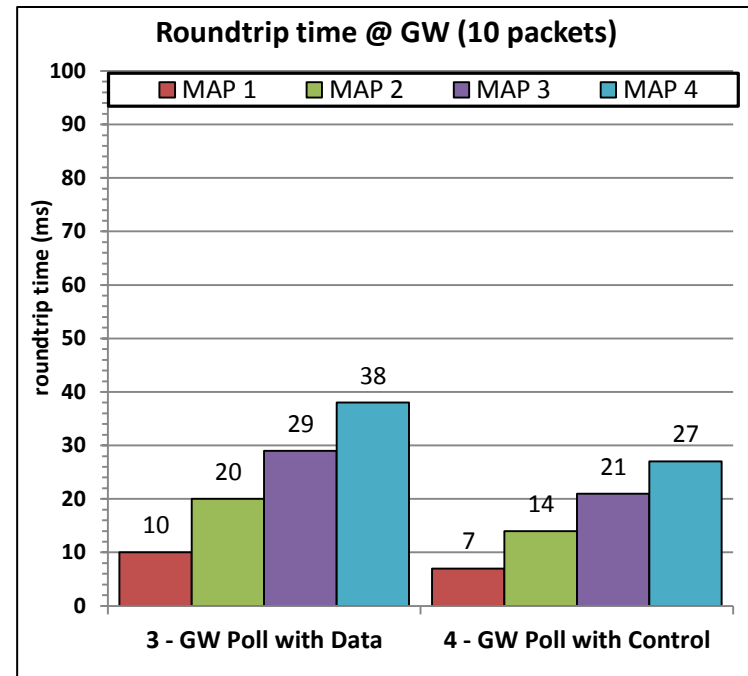
Example:



Case 3 data flow
Case 4 data flow

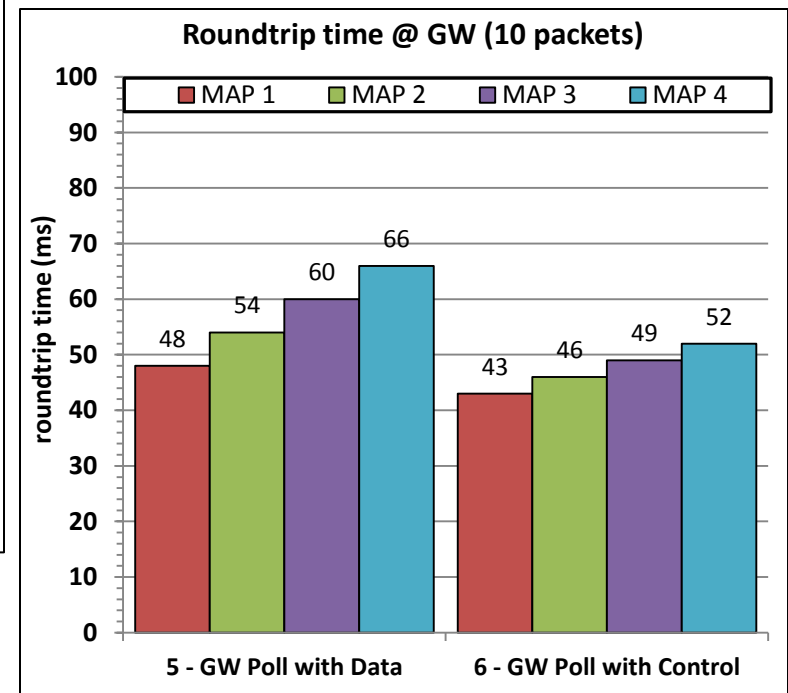
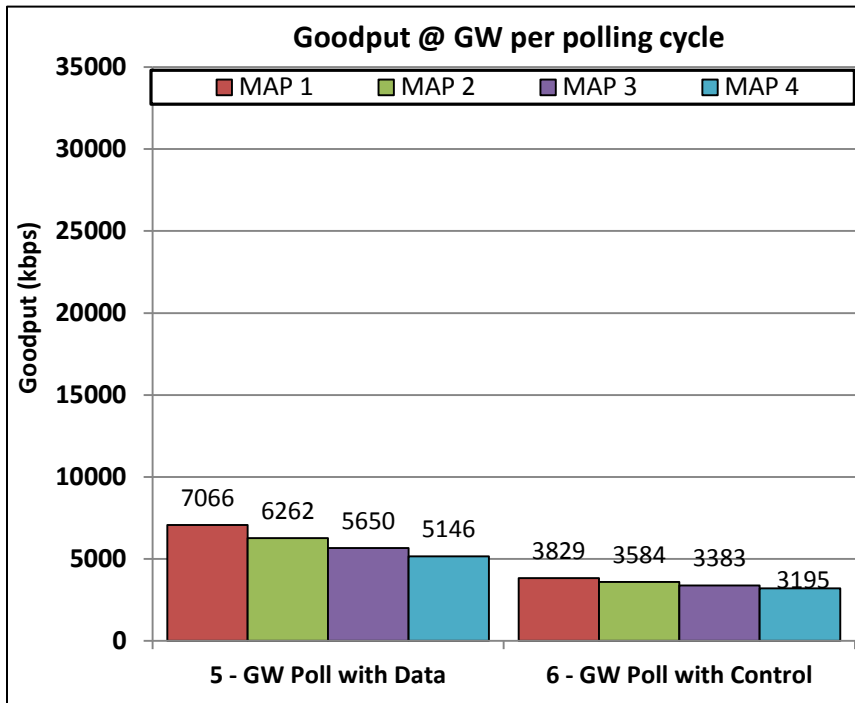
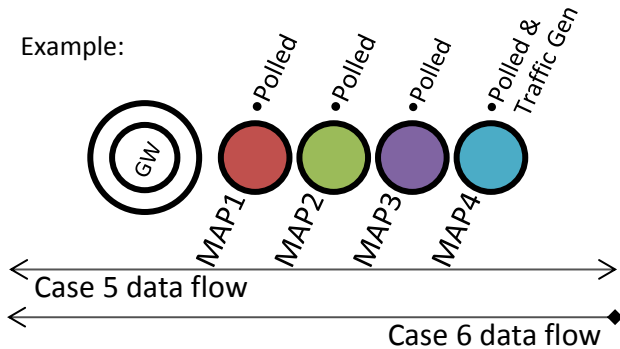


✓ No data packet losses



Results – Cases 5 and 6 | PACE | Full Polling | Only one active MAP

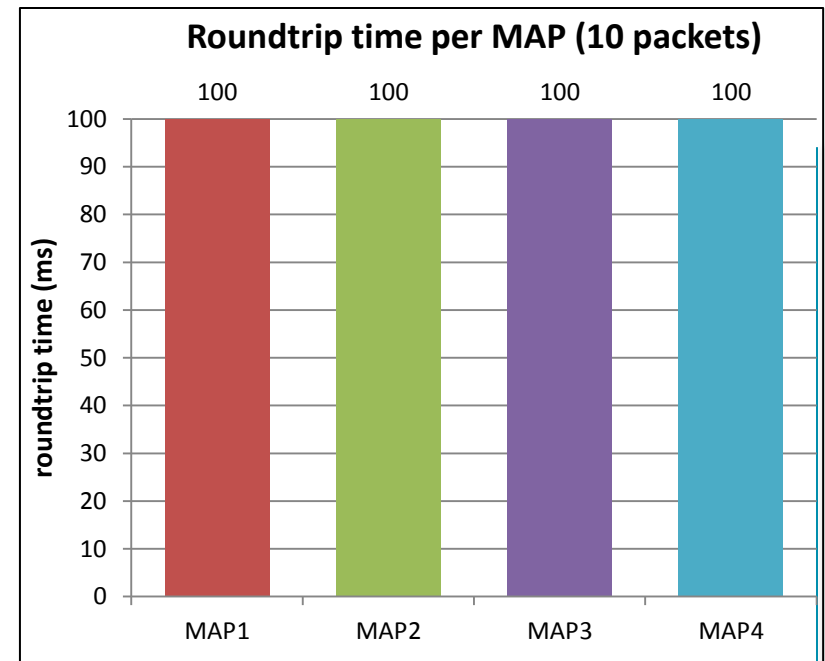
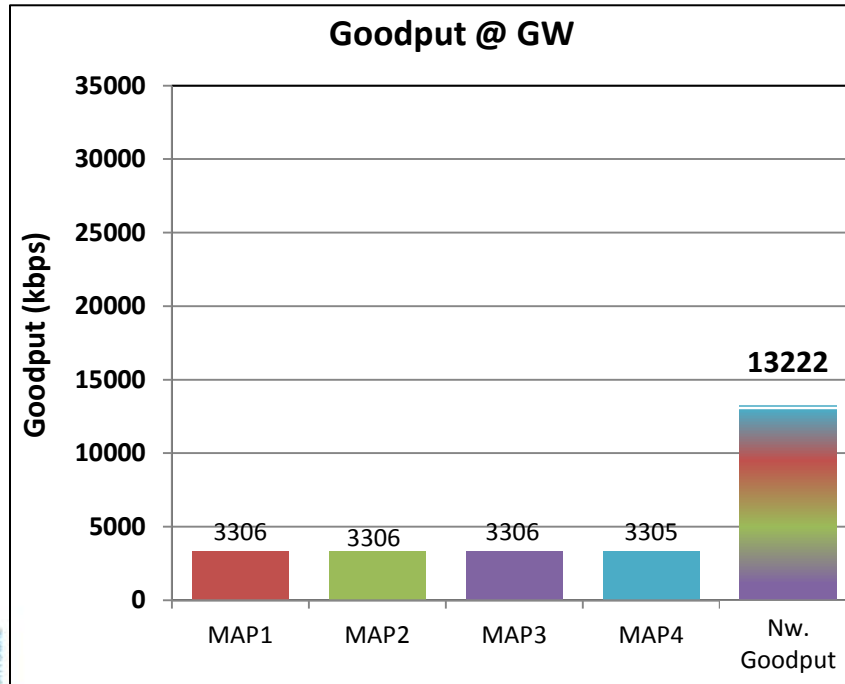
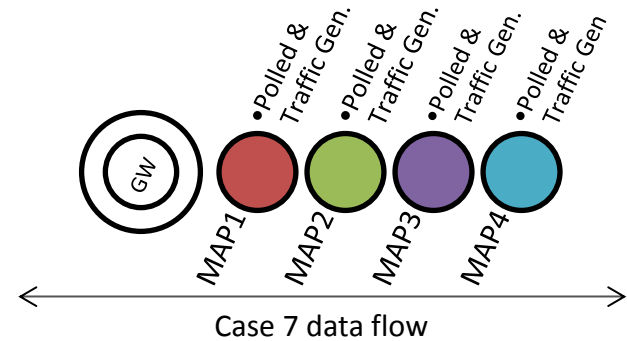
- Gateway polls all MAPs repeatedly
- Only one of the MAPs is active: Traffic generator
- In Case 5 the poll signal is implicit (Data)
- In Case 6 the poll signal is explicit (Control)



✓ No data packet losses

Results – Case 7 | PACE | Full Polling | All MAPs are active

- Gateway polls all MAPs repeatedly with Data
- All nodes are active : Traffic generators



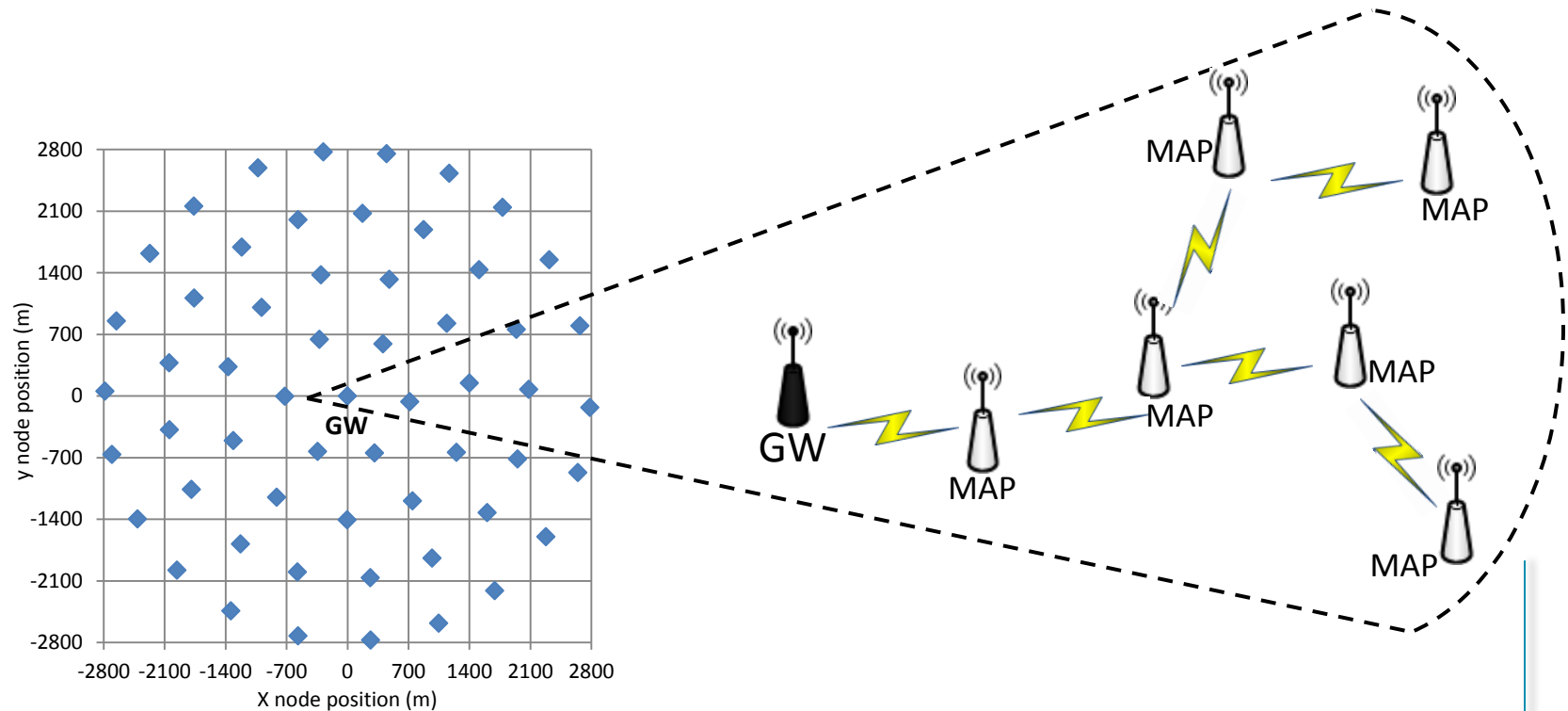
- ✓ No data packet losses
- ✓ Fairness at the data plane level

Results – Brief comparison

- Only MAP4 is active
 - CSMA/CA: 2,7 Mbps with 68 ms delay
 - PACE: at most 8.5 Mbps with 3.8 ms delay
- All MAPs are active
 - CSMA/CA at most in MAP4: 0.260 Mbps with 993ms delay
 - **Total network goodput: 1.1 Mbps**
 - PACE at MAP4: 3.3 Mbps with 10.0 ms delay
 - **Total network goodput : 13.2 Mbps**

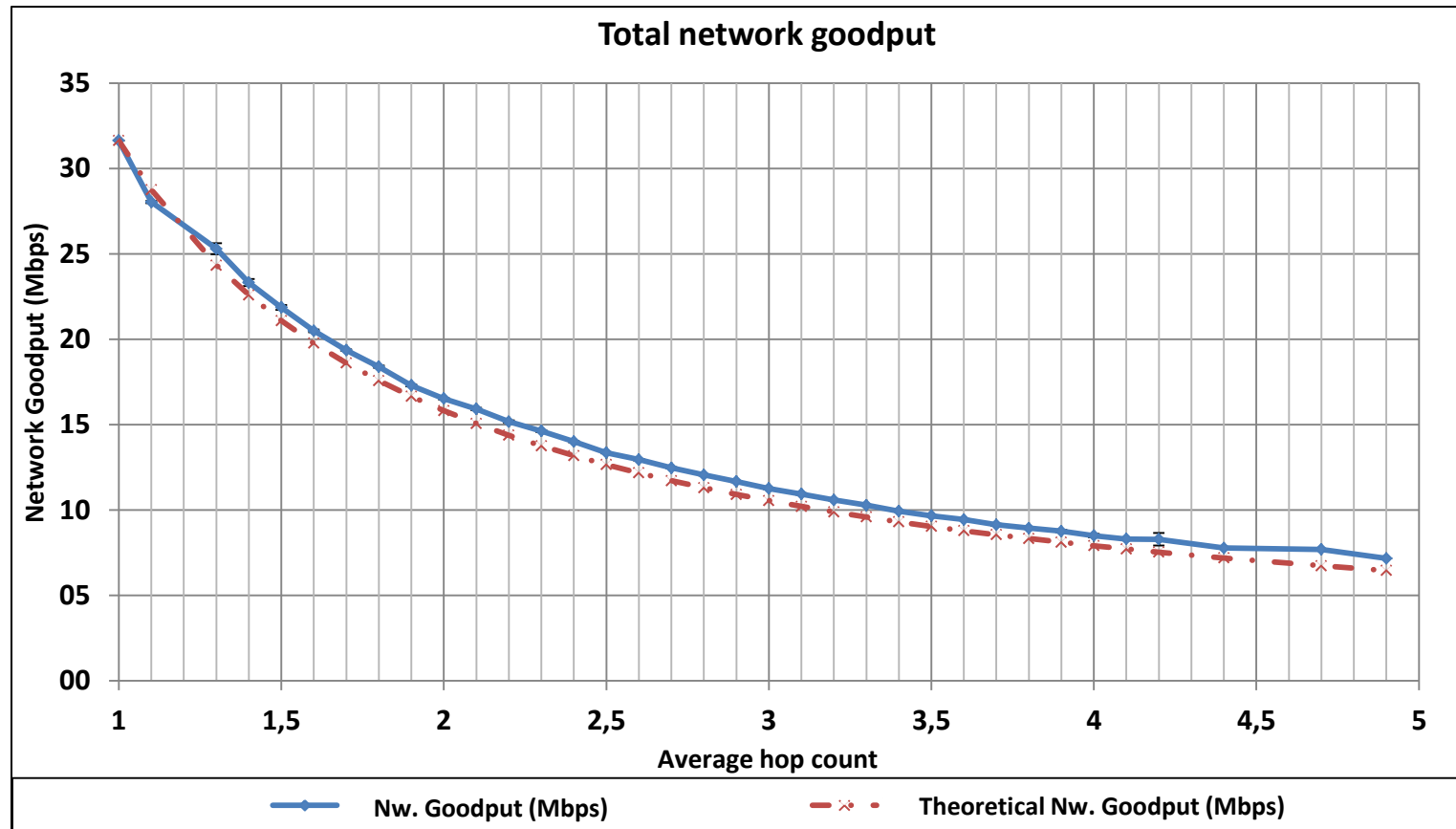
On average PACE is about 750% better!

Scenario 2 – Random Topology



- All MAPs are randomly placed on a radial arrangement
- Random logical topology
- Size of data frames: 2000 bytes
- RTS/CTS disabled
- Results are expressed as Goodput and Delay

Results – Case 8 | PACE | Full Polling | Random topology

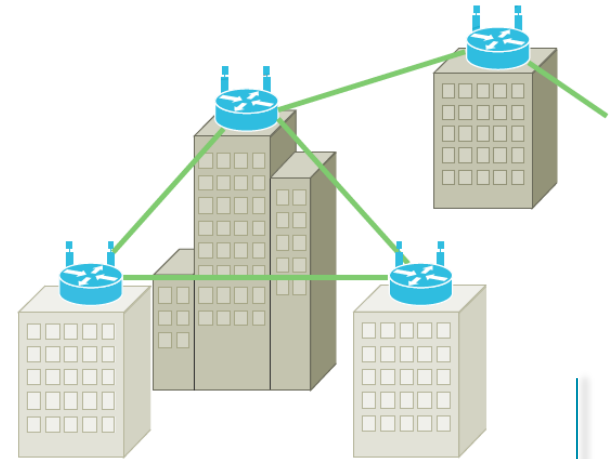


✓ Almost perfect match between theoretical and simulated values

$$C = C_{\text{Wi-Fi link}} / \text{AvgHopCount}$$

Conclusions and future work

- Routing, hidden nodes, congestion control and fairness are major challenges in Stub WMNs
- PACE is a simple and efficient approach
 - Fairness and predictable Stub WMN capacity
 - Easy and fast deployment
- Ongoing and future work
 - Temporal and spatial reuse
 - Integration with Smart Grid systems
 - Implementation in real 802.11-based testbeds
 - QoS: Load balancing, different priorities assigned to flows





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